## Reversible gates

- Landauer's Principle: Suppose a computer erases a single bit of information. The amount of energy dissipated into the environment is (the entropy of the environment increases by) at least $k_{B} T \ln 2$, where $k_{B}$ is Boltzmann's constant and $T$ is the temperature of the surrounding environment.
- Computers currently dissipate on the order of $500 k_{B} T \ln 2$ in energy for each elementary logical operation, so we are nowhere near the lower bound set by Landauer's Principle.
- Useful facts about XOR:

$$
\begin{aligned}
& z \oplus z=0 \\
& z \oplus 0=z \\
& z \oplus 1=\operatorname{NOT} z
\end{aligned}
$$

- AND, OR, NAND, NOR, XOR, and XNOR are not reversible.
- How to make XOR reversible? Add an extra output that just copies the $x$ input:

| x | y | x | XOR y |
| :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | 0 | 1 | 0 |
| 0 | 1 |  | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 |


| x | y | I | x | x XOR y |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | 0 | \| | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 0 |  |

- Reversible XOR is called the controlled-NOT or CNOT gate. Inputs: $x$ and $y$. Outputs: $x$ (unchanged) and $y^{\prime}=x \oplus y$. Input $x$ acts as a control bit. When $x=1, y$ is negated, since $1 \oplus y=$ NOT $y$. Otherwise $y$ is unchanged, since $0 \oplus y=y$.

CNOT : $|x, y\rangle \rightarrow|x, x \oplus y\rangle$

Matrix representation:
$C N O T=\left[\begin{array}{llll}1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0\end{array}\right]$

Gate symbol:


- $C N O T \star C N O T=I D E N$
- If we apply $C N O T$ to an input vector $V$, we can undo the effect by applying another $C N O T$ operation:
$C N O T \star C N O T \star V=I D E N \star V=V$
- Toffoli gate: controlled-controlled-NOT (CCNOT). Inputs: $x, y, z$. Outputs: $x^{\prime}, y^{\prime}, z^{\prime}$. Inputs $x$ and $y$ act as control bits, which are just copied to the outputs $x^{\prime}$ and $y^{\prime}$. When $x=1$ and $y=1, z$ is negated, since $z \oplus 1=$ NOT $z$. Otherwise $z$ is unchanged, since $z \oplus 0=z$.

- NOT from a Toffoli gate: set $x=1, y=1$ (just delete all rows from Toffoli truth table with $x=0$ or $y=0$ ):

| x | y | z | x | $y^{\prime} \quad z^{\prime}=$ NOT $z$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 |

- AND from a Toffoli gate: set $z=0$ (just delete rows with $z=1$ from Toffoli truth table):

- COPY from a Toffoli gate: set $x=1, z=0$ (just delete all rows from Toffoli truth table with $x=0$ or $z=1$ ). Input $y$ gets copied to outputs $y^{\prime}$ and $z^{\prime}$ :

| x | y | z | x' | $y^{\prime}$ | z' |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 |

- Fredkin gate: controlled-SWAP
$|0, y, z\rangle \rightarrow|0, y, z\rangle$
$|1, y, z\rangle \rightarrow|1, z, y\rangle$
When $x=1$, the inputs $y$ and $z$ are swapped. Otherwise they are unchanged. Truth table:

- AND from Fredkin: set $y=0$ (delete all rows with $y=1$ ). Output $y^{\prime}=x$ AND $z$

| x | y | z | $\mid$ | x | $\mathrm{y}^{\prime}=\mathrm{x}$ AND z | $\mathrm{z}^{\prime}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| -0 | 0 | 0 | $\mid$ | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 0 | $\mid$ | 1 | 0 | 0 |
| 1 | 0 | 1 | $\mid$ | 1 | 1 | 0 |

- NOT and COPY from Fredkin: set $y=1, z=0$. Output $y^{\prime}=$ NOT $x$. Input $x$ gets copied to outputs $x$ and $z^{\prime}$ :

- Since we can make NOT and AND, we can make NAND, so we can make anything from Fredkin gates.

